



Relevance of Backward Error

What do we gain from a bound on backward error like

$$rac{\|oldsymbol{x}-oldsymbol{\hat{x}}\|}{\|oldsymbol{x}\|} \leq \epsilon?$$



Demo: Backward Stability by Example [cleared]

Getting into Trouble with Accuracy and Stability

How can I produce inaccurate results?

Apply an inaccurate method

- Apply an unstable method to a well-conditioned problem
- > Apply any type of method to an ill-conditioned problem

Wanted: Real Numbers... in a computer

Computers can represent integers, using bits:

$$23 = 1 \cdot 2^4 + 0 \cdot 2^3 + 1 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 = (10111)_2$$

How would we represent fractions?

$$23.625 = (10111.000) = 25.625 = (10111.000) = 25.625 = (10110.000) = 25.625 = 25.625$$

Fixed-Point Numbers

Suppose we use units of 64 bits, with 32 bits for exponents ≥ 0 and 32 bits for exponents < 0. What numbers can we represent?



How many 'digits' of relative accuracy (think relative rounding error) are available for the smallest vs. the largest number?



Floating Point Numbers



Floating Point: Implementation, Normalization

Previously: Consider *mathematical* view of FP. (via example: (1101)₂) Next: Consider *implementation* of FP in hardware.

Do you notice a source of inefficiency in our number representation?

Implementing Arithmetic

How is floating point addition implemented? Consider adding $a = (1.101)_2 \cdot 2^1$ and $b = (1.001)_2 \cdot 2^{-1}$ in a system with three stored bits (four total) in the significand.



Implementing Arithmetic

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Rough algorithm:

- 1. Bring both numbers onto a common exponent
- 2. Do grade-school addition from the front, until you run out of digits in your system.
- 3. Round result.

 $a = 1. \quad 101 \cdot 2^{1}$ $b = 0. \quad 01001 \cdot 2^{1}$ $a + b \approx 1. \quad 111 \cdot 2^{1}$

Unrepresentable numbers?

Can you think of a somewhat central number that we cannot represent as

$$x = (1. _ _ _ _ _ _)_2 \cdot 2^{-p}?$$

Demo: Picking apart a floating point number [cleared]

Subnormal Numbers

What is the smallest representable number in an FP system with 4 stored bits (5 total) in the significand and a stored exponent range of [-7, 8]?

_ _ _ _ _ _ _ _ _ (done v/a demo)

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Subnormal Numbers, Attempt 2

What is the smallest representable number in an FP system with 4 stored bits in the significand and a (stored) exponent range of [-7, 8]?

	(donc	V/a	demo)		
Why learn about subnormals?					

Subnormal Numbers, Attempt 2

What is the smallest representable number in an FP system with 4 stored bits in the significand and a (stored) exponent range of [-7, 8]?

- Can go way smaller using the special exponent (turns off the leading one)
- Assume that the special exponent is -7; interpreted as -6.
- So: $(0.0001)_2 \cdot 2^{-6}$ (with four digits after the point stored).
- Numbers with the special exponent are called *subnormal* (or *denormal*) FP numbers. Technically, zero is also a subnormal.

Why learn about subnormals?

Subnormal FP is often slow: not implemented in hardware.

Many compilers support options to 'flush subnormals to zero'.