

CS 450: Numerical Analysis

Lecture 5

Chapter 2 – Linear Systems

Solving Linear Systems

Edgar Solomonik

Department of Computer Science
University of Illinois at Urbana-Champaign

January 31, 2018

Solving Triangular Systems

- ▶ $Lx = b$ if L is lower-triangular is solved by forward substitution:

$$\begin{array}{rcl} l_{11}x_1 = b_1 & & x_1 = b_1/l_{11} \\ l_{21}x_1 + l_{22}x_2 = b_2 & \Rightarrow & x_2 = (b_2 - l_{21}x_1)/l_{22} \\ l_{31}x_1 + l_{32}x_2 + l_{33}x_3 = b_3 & & x_3 = (b_3 - l_{31}x_1 - l_{32}x_2)/l_{33} \\ & & \vdots \\ & & \vdots \end{array}$$

- ▶ Computational complexity of forward/backward substitution:

$$T(n) = O(n^2)$$

$$\frac{n^2}{2} \text{ mults} \quad \frac{n^2}{2} \text{ additions}$$

Solving Triangular Systems

- ▶ Existence of solution to $Lx = b$:

May not exist if some $L_{ii} = 0$, since then L is singular

- ▶ Invertibility of L and existence of solution:

sol'n may exist even if $L_{ii} = 0$, but is not unique

L is full rank if each $L_{ii} \neq 0$
(determinant shows this formally, since $\det(L) = \prod_i L_{ii}$,

but can also see that $Lx = b$ when b is not in the span of the last $n-1$ cols, since $L_{ii} \neq 0$)

Properties of Triangular Matrices

- ▶ $XY = Z$ is lower triangular if X and Y are both lower triangular:

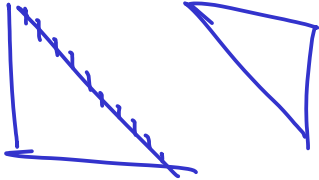
$$\triangle \triangle \approx \triangle \quad \nabla \nabla = \nabla$$

- ▶ L^{-1} is lower triangular if it exists:

$$\triangle^{-1} = \triangle$$

LU Factorization

- ▶ An **LU factorization** consists of a unit-lower-triangular **factor** L and upper-triangular factor U such that $A = LU$:



n^2 variables in $L, U \leftarrow \frac{n(n-1)}{2}$
 n^2 variables in A

Gaussian Elimination

- ▶ The LU factorization may not exist:

Consider matrix $\begin{bmatrix} 1 & 2 \\ 2 & 4 \\ 0 & 3 \end{bmatrix}$

$$= \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} \begin{bmatrix} u_{11} & u_{12} \\ 0 & u_{22} \\ 0 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 2 & 0 \\ 2 & 1 & 0 \\ 0 & l_{32} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 2 & u_{22} \\ 0 & 0 \end{bmatrix}$$

Contradiction

$$\begin{aligned} & \frac{3}{l_{22}} \quad u_{22} = 0 \\ & \uparrow l_{22} \\ & l_{21} l_{22} = 3 \\ & 2 l_{21} = 0 \\ & \downarrow \\ & 0 \end{aligned}$$

- ▶ Permutation of variables enables us to transform the linear system so the LU factorization does exist:

Gaussian Elimination Algorithm

- ▶ Algorithm for factorization is derived from equations given by $A = LU$:

$$\begin{cases} \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} = \begin{bmatrix} L_{11} & \\ L_{21} & L_{22} \end{bmatrix} \begin{bmatrix} U_{11} & U_{12} \\ & U_{22} \end{bmatrix} \end{cases}$$

$$\begin{array}{l}
 \begin{array}{l}
 A_{11} = L_{11} U_{11} \quad L_{11} U_{12} = A_{12} \quad \triangle \square = \square \\
 L_{21} U_{11} = A_{21} \quad \square \nabla = \square \rightarrow \nabla = \square \quad \left\{ \begin{array}{l} L_{21} U_{12} + L_{22} U_{22} = A_{22} \end{array} \right.
 \end{array}
 \end{array}$$

- ▶ The k th column of L is given by the k th **elementary matrix** M_k :

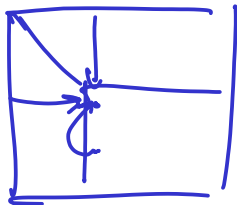
$$\begin{array}{l}
 \underbrace{A_{22} - L_{21} U_{12}}_{\text{Solve complete}} \\
 L_{21} U_{12} + L_{22} U_{22} = A_{22} \\
 A_{22} - L_{21} U_{12} = L_{22} U_{22}
 \end{array}$$

Elimination Matrices

- ▶ **An elimination matrix M_k satisfies the following properties:**

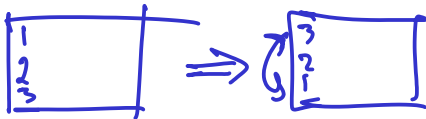
Gaussian Elimination with Partial Pivoting

- ▶ **Partial pivoting** permutes rows to make divisor u_{jj} maximal at each step:



$$A = PLU$$

$$P^T A = LU$$



- ▶ A row permutation corresponds to an application of a **row permutation matrix** $P_{jk} = I - (e_j - e_k)(e_j - e_k)^T$:

$$P^T P = I$$

$$\begin{bmatrix} 1 \\ -1 \\ -1 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ & 1 \end{bmatrix} = \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix}$$

$$P_{jk} = \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & 1 & \\ & & & \ddots \end{bmatrix}$$

Complete Pivoting and Error Bounds

- ▶ **Complete pivoting** permutes rows and columns to make divisor u_{ii} is maximal at each step:

partial pivoting $|L|_{1,j} \leq 1$

- ▶ For LU, the backward error δA , so that $\hat{L}\hat{U} = A + \delta A$, satisfies bound $|\delta A_{ij}| \leq \epsilon (|\hat{L}| \cdot |\hat{U}|)_{ij}$:

