

# CS 450: Numerical Analysis

Lecture 20

Chapter 7 Interpolation

Chebyshev Interpolation

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# Orthogonal Polynomials

- ▶ Recall that good conditioning for interpolation is achieved by constructing a well-conditioned Vandermonde matrix, which is the case when the columns (corresponding to each basis function) are orthonormal. To construct robust basis sets, we introduce a notion of *orthonormal functions*:

$$\langle f, g \rangle_w = \int_{-\infty}^{\infty} f(x) g(x) dx$$

$$\|f\| = \sqrt{\langle f, f \rangle_w} \quad \leftarrow \text{normalization (size measure)}$$

Legendre normalization  $\rightarrow$  by  $f(1)$   
end of interval  $\uparrow$

## Legendre Polynomials

- ▶ The Gram-Schmidt orthogonalization procedure can be used to obtain an orthonormal basis with the same span as any given arbitrary basis:

$$\left[ \begin{array}{cccc} 1 & x & x^2 & x^3 \\ \downarrow & \downarrow & & \\ f_1(x) & f_2(x) & \dots & \end{array} \right] \quad e_i(x) = f_i(x) - \sum_{j=1}^{i-1} \langle f_j(x), e_i(x) \rangle \cdot e_j(x)$$

- ▶ The Legendre polynomials are obtained by Gram-Schmidt on the monomial basis, with normalization done so  $\hat{\phi}_i(1) = 1$  and  $w(t) = \begin{cases} 1 : -1 \leq t \leq 1 \\ 0 : \text{otherwise} \end{cases}$

## Basis Orthogonality and Conditioning

- ▶ To obtain perfectly conditioned Vandermonde system, want inner products of different columns to be zero:

$A$  is orthogonal!

$$A^T A = I$$

columns of  $A$ ,  $a_1, \dots, a_n$

$$\langle a_i, a_j \rangle = \sum_{k=1}^n w(t_k) e_i(t_k) e_j(t_k) = \delta_{ij}$$

- ▶ These inner products should be close to zero, if they are a suitable quadrature rule for our weighted functional inner product:

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$$\int_{-\infty}^{\infty} w(t) e_i(t) e_j(t) dt \approx \sum_{k=1}^n w(t_k) e_i(t_k) e_j(t_k)$$

# Chebyshev Basis

- ▶ **Chebyshev polynomials**  $\phi_j(t) = \cos(j \arccos(t))$  and **Chebyshev nodes**  $t_i = \cos\left(\frac{2i-1}{2n}\pi\right)$  provide a way to pick **nodes**  $t_1, \dots, t_n$  along with a basis, to yield perfect conditioning:

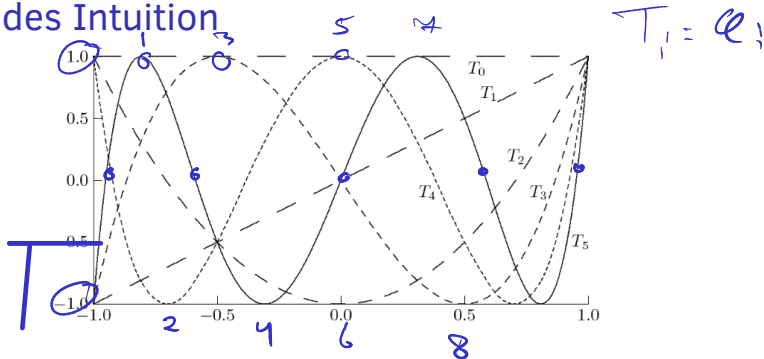
$$\phi_j(t_i) = \cos\left(j \frac{2i-1}{2n} \pi\right) = V_{ji} \quad \left| \begin{array}{l} \text{orthogonal polynomial} \\ \text{w.r.t. } w(t) = 1/(1-t^2) \end{array} \right.$$



$$\begin{aligned} & \sum_{k=1}^n \phi_i(t_k) \phi_j(t_k) \\ &= \sum_{k=1}^n \cos\left(i \frac{2k-1}{2n} \pi\right) \cos\left(j \frac{2k-1}{2n} \pi\right) \\ &= 0 \quad \text{if } i \neq j \end{aligned}$$

$$\sum_{k=1}^n V_{ik} V_{kj} = 0 \quad \text{if } i \neq j \quad V^T V = 0$$

# Chebyshev Nodes Intuition



- ▶ Note equi-alteration property, successive extrema of  $T_k$  have opposite sign:

and equal magnitude

- ▶ Set of  $k$  Chebyshev nodes are given by zeros of  $T_k$

$n$ -nodes get zeros  $T_n$   
 use basis functions  $T_0 \dots T_{n-1}$   $n = 1, 2, \dots$

## Orthogonal Polynomials and Recurrences

- ▶ The Newton polynomials could be obtained by a two-term recurrence
- ▶ Legendre and Chebyshev polynomials also satisfy three-term recurrence, for Chebyshev

$$\phi_{i+1}(t) = 2t\phi_i(t) - \phi_{i-1}(t)$$

$$\begin{aligned} \varphi_0(t) &= 1 \\ \varphi_1(t) &= x \end{aligned}$$

- ▶ In fact all orthogonal polynomials satisfy some recurrence of the form,

$$\varphi_{i+1}(t) = (\alpha_i + \beta_i t) \varphi_i(t) - \gamma_i \varphi_{i-1}(t)$$

# Error in Interpolation

True function

Given degree  $n$  polynomial interpolant  $\tilde{f}$  of  $f$  induction on  $n$  shows that  $E(x) = f(x) - \tilde{f}(x)$  has  $n$  zeros  $x_1, \dots, x_n$  and there exist  $y_1, \dots, y_n$  such that

$[x_1, x_2] \dots [x_n, x_n]$

interpolant  $\nearrow$

$$E(x) = \int_{x_1}^x \int_{y_1}^{w_0} \dots \int_{y_n}^{w_{n-1}} \underbrace{f^{(n+1)}(w_n)}_{\text{degree}} dw_n \dots dw_0 \quad (1)$$

inductive hypothesis  $\swarrow$

$$E(x) = \underbrace{E(x_i)}_0 + \int_{x_i}^x \underbrace{E'(x)}_{\text{degree}} dx$$

$$E(x_i) = 0$$

$$E(x_i) - E(x_{i-1}) = \int_{x_{i-1}}^{x_i} E'(x) dx = 0$$

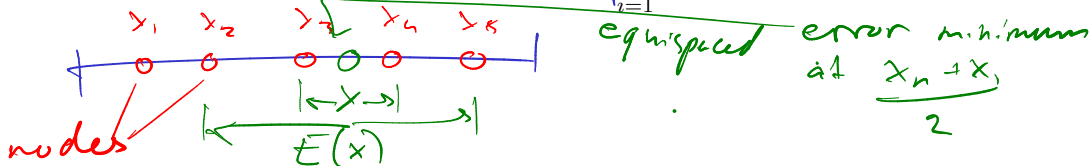
$\exists y_i \in [x_{i-1}, x_i]$   
 $E'(y_i) = 0$



# Interpolation Error Bounds

- ▶ The error bound implies that

$$|E(x)| \leq \frac{\max_{s \in [x_1, x_n]} |f^{(n+1)}(s)|}{n!} \left| \prod_{i=1}^n (x - x_i) \right| \quad \text{for } x \in [x_1, x_n]$$



- ▶ Letting  $h = x_n - x_1$  (often also achieve same for  $h$  as the node-spacing  $x_{i+1} - x_i$ ), we obtain

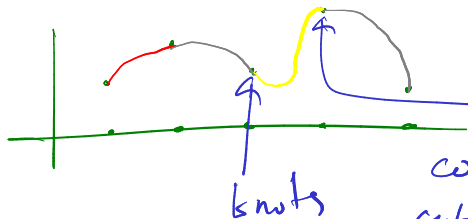
$$|E(x)| \leq \frac{\max_{s \in [x_1, x_n]} |f^{(n+1)}(s)|}{n!} h^n = O(h^n) \quad \text{for } x \in [x_1, x_n]$$

more nodes  $\rightarrow$  increase  $n$

decreasing  $h \rightarrow$   $n$ th order convergence  
 motivates piecewise polynomial interp.

# Piecewise Polynomial Interpolation

- ▶ The  $k$ th piece of the interpolant is a polynomial in  $[x_i, x_{i+1}]$



enforce differentiability  
(once)  
common to use  
cubic interpolants for each piece

- ▶ **Hermite** interpolation ensures consecutive interpolant pieces have same derivative at each **knot**  $x_i$ :

$f_i(x)$  is  $i$ th piece defined  $(x_i, x_{i+1})$

$$\frac{df_i}{dx}(x_{i+1}) = \frac{df_{i+1}}{dx}(x_{i+1})$$

# Spline Interpolation

- ▶ A **spline** is a  $(k - 1)$ -time differentiable piecewise polynomial of degree  $k$ :

whole interpolant  
 cubic spline has  $4(n-1)$  parameters  
 natural spline sets 2nd derivative to 0 at  $x_1$  and  $x_n$   
 $k+1$  pieces/nodes

$2(n-1)$  equations for interpolants,  $n-1$  for each deriv.

- ▶ The resulting interpolant coefficients are again determined by an appropriate generalized Vandermonde system:

$$f_1(x) = \alpha_0 + \alpha_1 x + \alpha_2 x^2 + \alpha_3 x^3$$

$$f_2(x) = \beta_0 + \beta_1 x + \dots$$

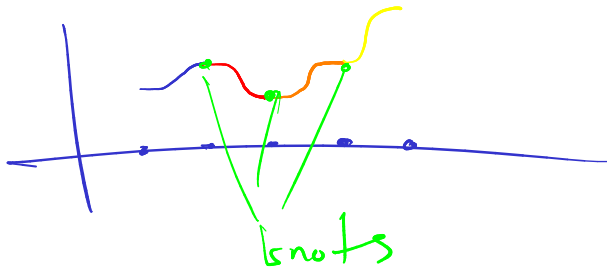
$$Vc = y$$

$$\begin{bmatrix} \vdots \\ \vdots \\ \vdots \\ \vdots \end{bmatrix} \begin{bmatrix} \vdots \\ \vdots \\ \vdots \\ \vdots \end{bmatrix} = \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \alpha_2 \\ \alpha_3 \\ \vdots \\ \vdots \\ \beta_0 \\ \beta_1 \\ \vdots \end{bmatrix}$$

$2(n-1)$   
 $+ 1-2$   
 $+ n-2$   
 $4n - 6$

# Lecture 21: Recap Interpolation

## Piecewise interpolation



at each knot - both pieces interpolate  
- smoothness, differentiability

# Vandermonde Systems




spline interpolant: degree  $k$ ,  $(k-1)$ -times differentiable

natural spline - 2nd derivative <sup>e.g.  $k=3$</sup>  at  $x_1, x_n$  is zero

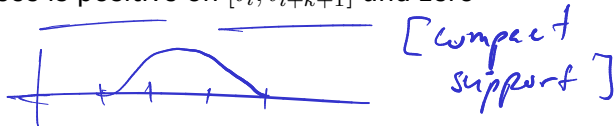
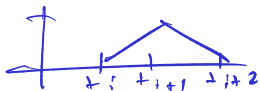
# B-Splines

*B-splines* provide an effective way of constructing splines from a basis:

- ▶ The basis functions can be defined recursively with respect to degree.

$$e_i^k = v_i^k e_i^{k-1} + (1 - v_{i+1}^k) e_{i+1}^{k-1} \quad \left| \quad e_i^0(t) = \begin{cases} 1 & \text{if } t \in [t_i, t_{i+1}] \\ 0 & \text{otherwise} \end{cases}$$
$$v_i^k(t) = \frac{t - t_i}{t_{i+k} - t_i}$$


- ▶ The  $i$ th degree  $k$  polynomial piece is positive on  $[t_i, t_{i+k+1}]$  and zero everywhere else



- ▶ All possible splines of degree  $k$  with notes  $\{t_i\}_{i=1}^n$  can be represented in the basis.

so coefficients of neighboring splines are dependent  $\rightarrow$  banded Vandermonde system